

Using Minimal Arrival Curves to Derive Per-Flow Performance Bounds in (Tandem) Networks with Complex Feedback Structures ¹

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Outline

- 1 Previous NC Work on Feedback
- 2 Extending Analysis to Interdependent Feedback Constraints
- 3 Deriving Per-Flow Performance Bounds
- 4 Conclusion

Previous NC Work on Feedback

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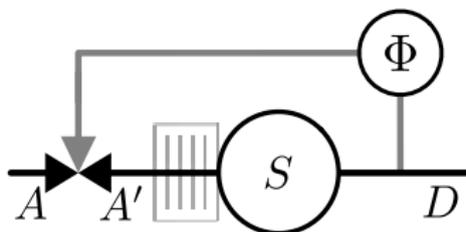


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- System with some sort of signaling mechanism (finite buffer, ...)

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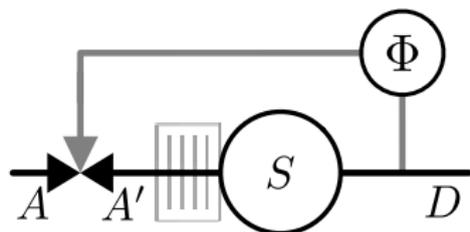


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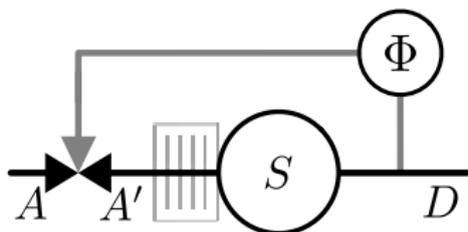


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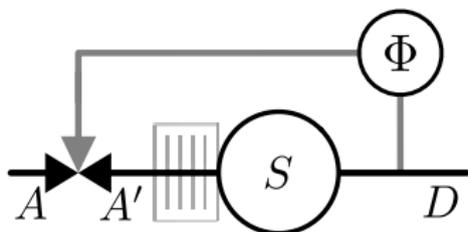


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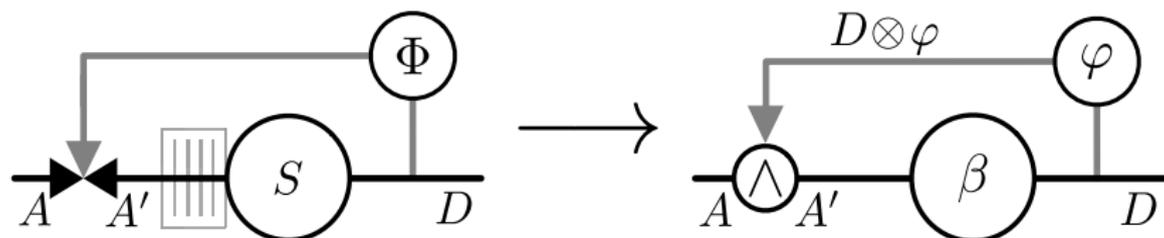


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- Signaling is modeled as backwards loop that feeds back into dataflow
- Feedback is modeled by some function Φ
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- Call φ with $D \circ \Phi \geq D \otimes \varphi$ a feedback curve

What Do We Want To Analyze?

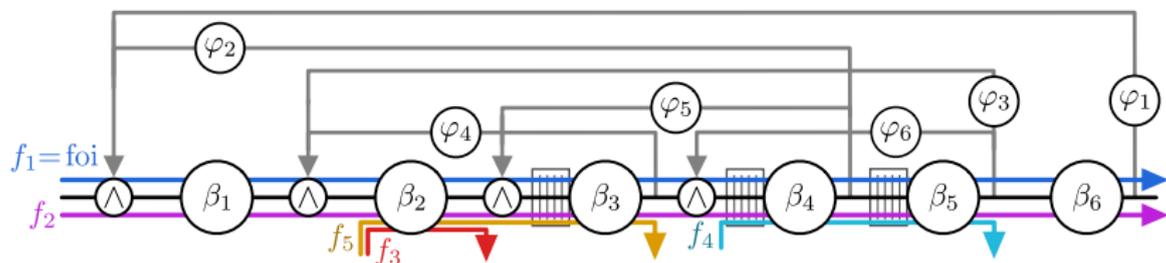


Figure: Example network with various feedback constraints and flows.

- Network with a multitude of feedback constraints that potentially have interdependencies
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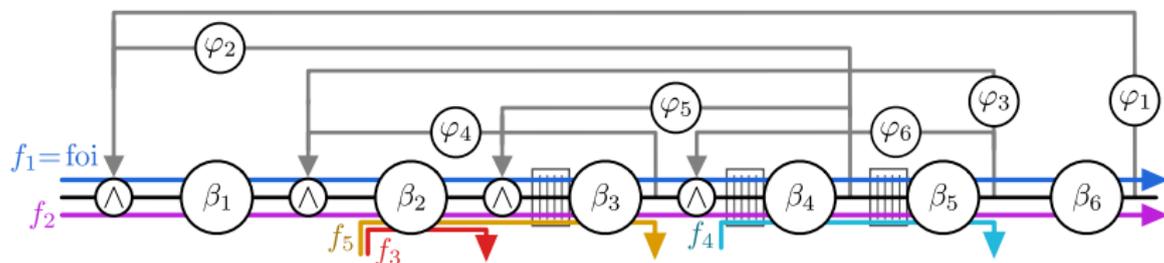


Figure: Example network with various feedback constraints and flows.

- Network with a multitude of feedback constraints that potentially have interdependencies
- Arbitrary number of interfering flows
- Goal: Obtain per-flow performance bounds for any scheduling policy
 - Not just FIFO...

Closed-Loop System

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- Issue: not min-plus linear [Boudec and Thiran, 2001], hence cannot be (directly) analyzed using NC

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- Issue: not min-plus linear [Boudec and Thiran, 2001], hence cannot be (directly) analyzed using NC
- Solution: transformation into *open-loop system* [Chang, 2000]

Theorem

Let an arrival process A traverse a system S offering service curve $\beta \in \mathcal{F}$. The system is subject to a feedback curve φ . If

$$D \geq (A \wedge (D \otimes \varphi)) \otimes \beta,$$

then

$$D \geq A \otimes \beta \otimes (\beta \otimes \varphi)^*.$$

Transformation Into Open-Loop System

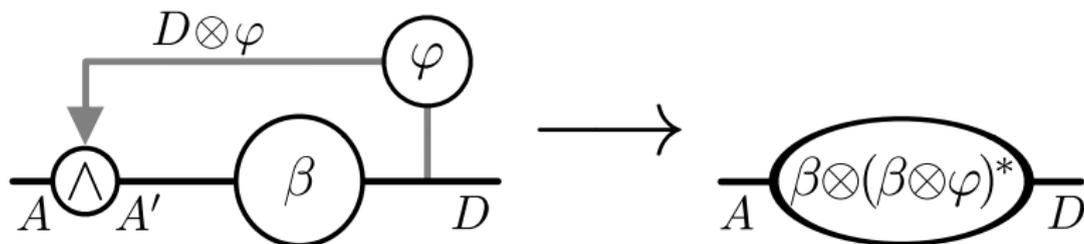


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- Using the Theorem, we obtain an open-loop system

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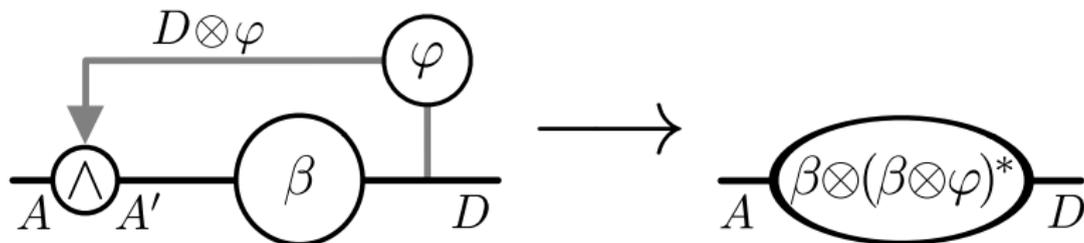


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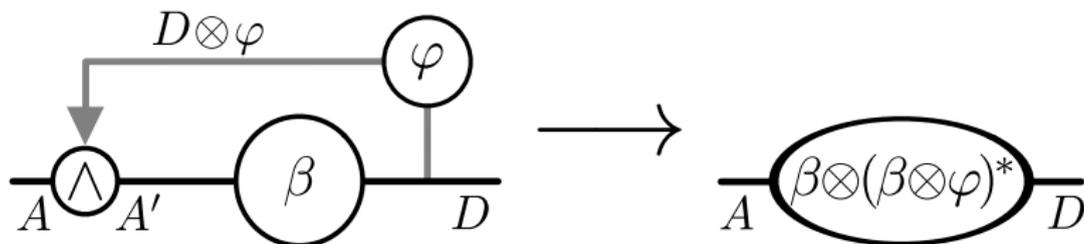


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- Using the Theorem, we obtain an open-loop system
- Feedback is still captured in the system, but encoded into the service curve directly
- Resulting system is min-plus linear \rightarrow can be analyzed with normal NC methods

Generalization To Arbitrary Length

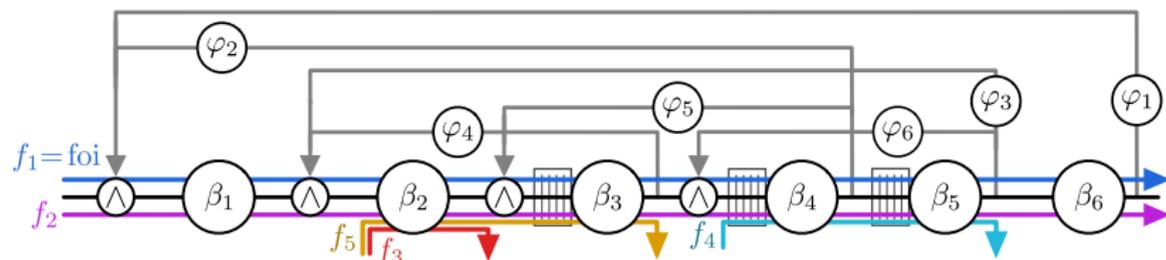


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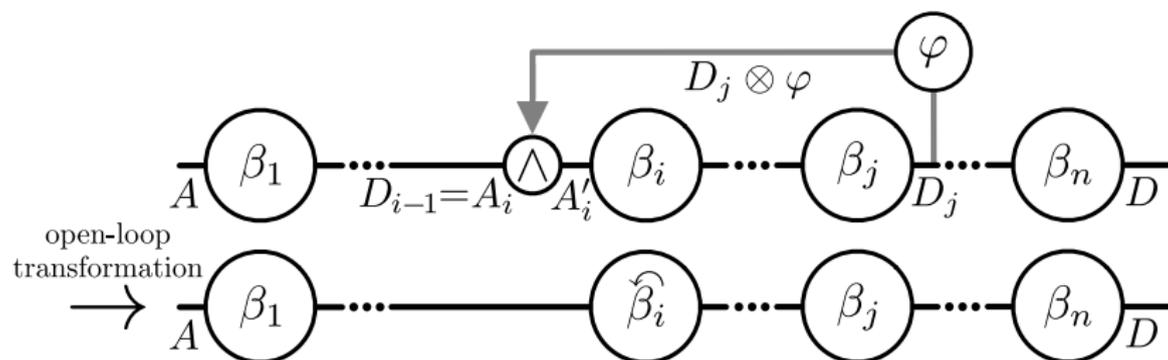


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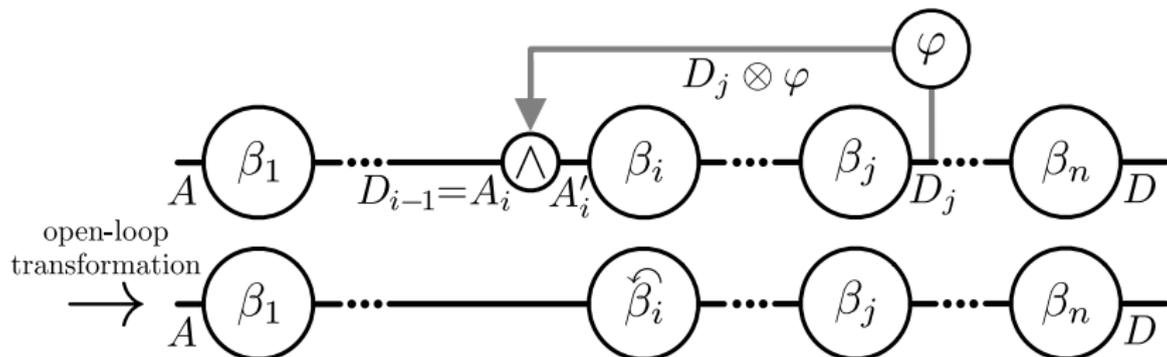


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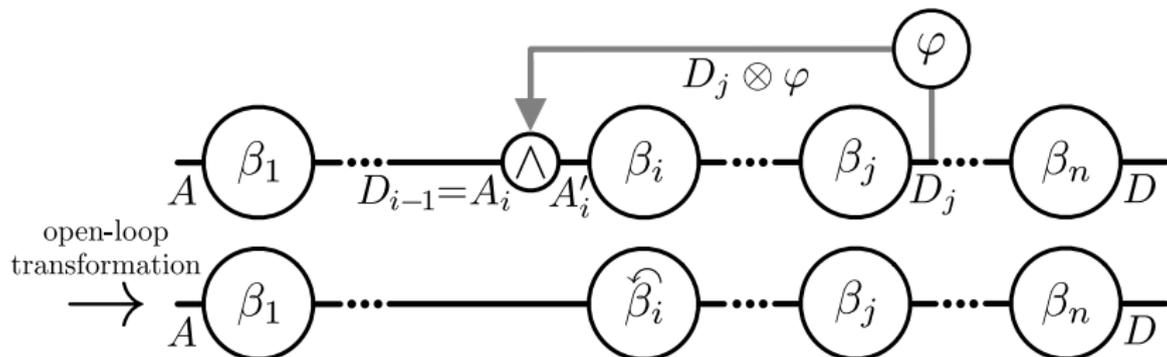


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- Call this Proposition for remainder of the talk

Extending Analysis to Interdependent Feedback Constraints

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- Some existing results on overlapping interdependency [Chang, 2000, Bose et al., 2006, Bouillard et al., 2009]

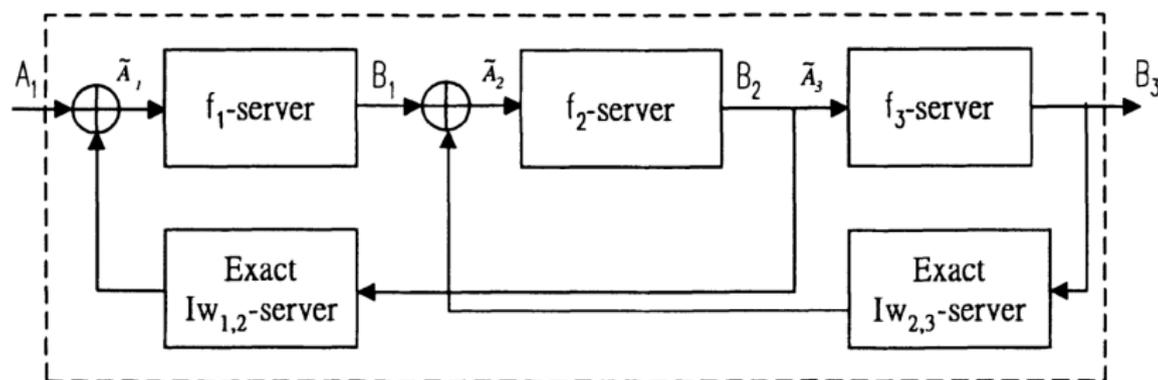


Figure: Overlapping interdependency [Chang, 2000].

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- What about having multiple constraints in a single system?
- What if the constraints have an interdependency?
- Some existing results on overlapping interdependency [Chang, 2000, Bose et al., 2006, Bouillard et al., 2009]
- But, not the only interdependence type!

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- Establish lexicographical order over all existing feedback constraints $F_l = (i_l, j_l, \varphi_l)$, $l = 1, \dots, k$, in the system as follows

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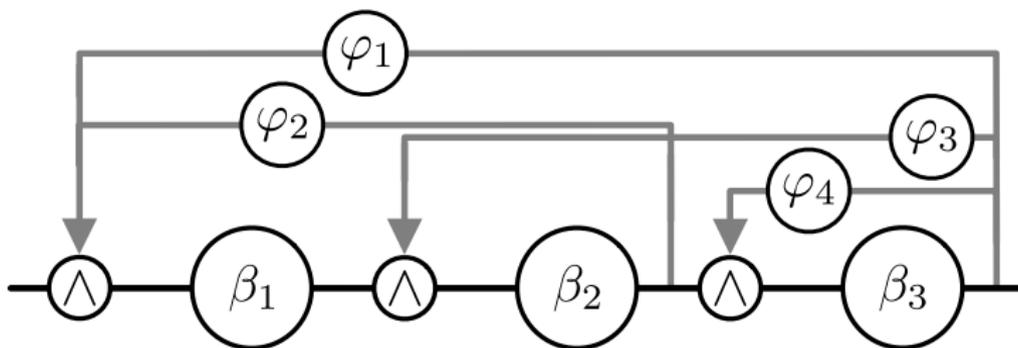


Figure: Network with interdependent feedback constraints.

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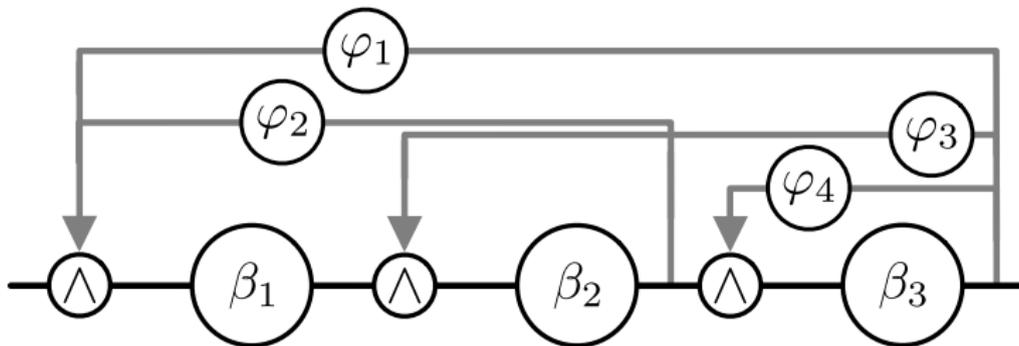


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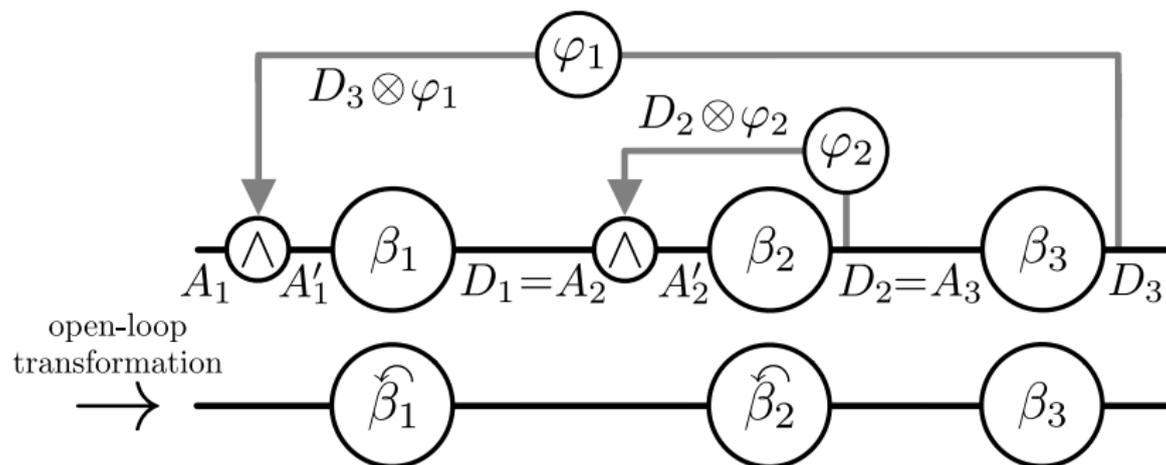


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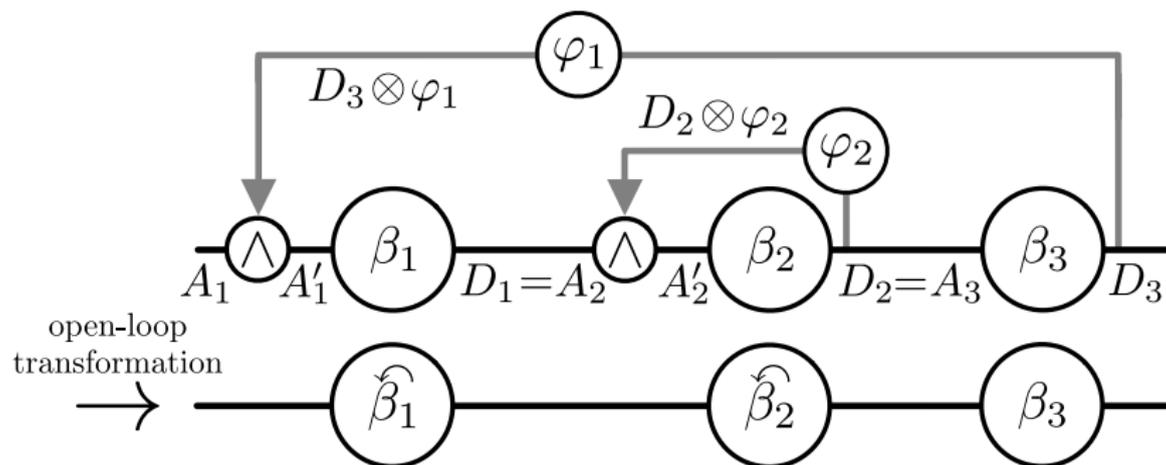


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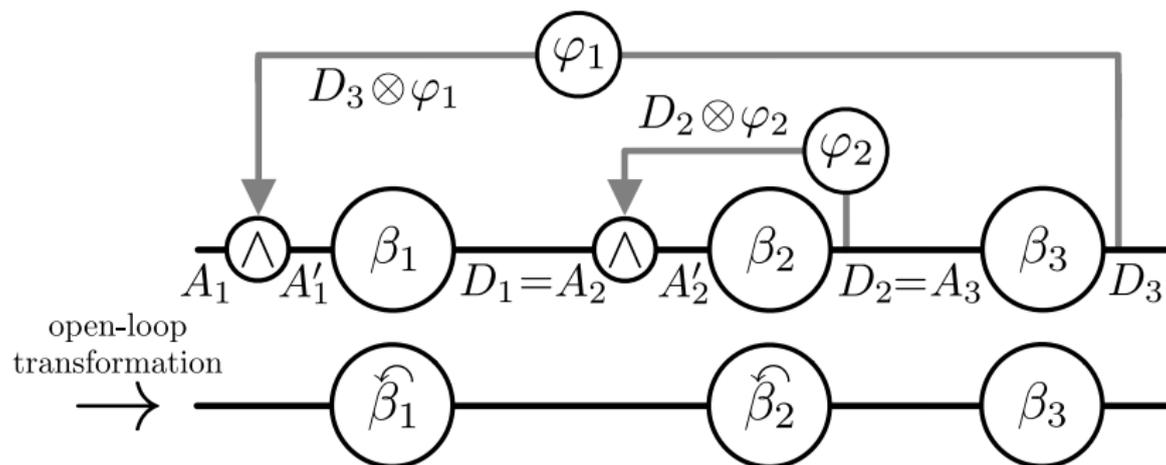


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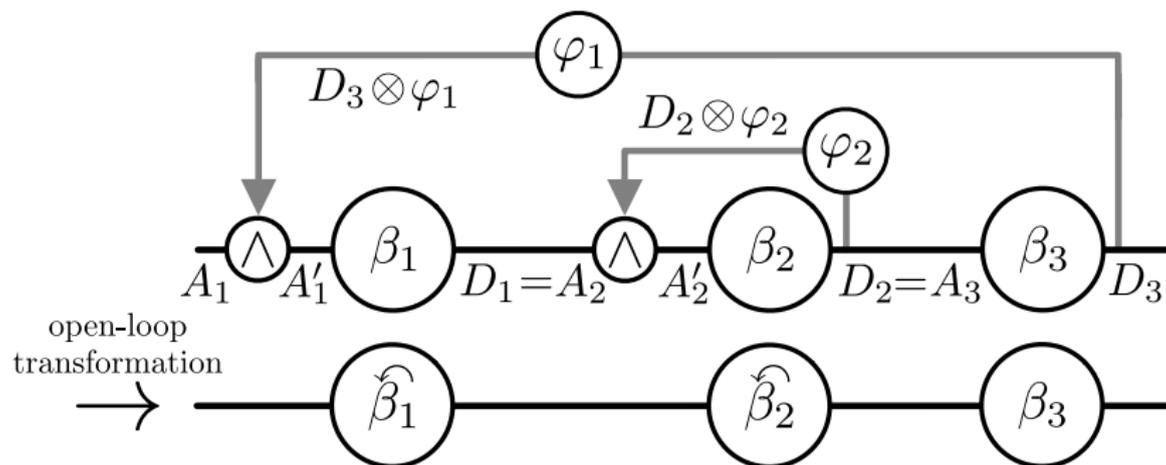


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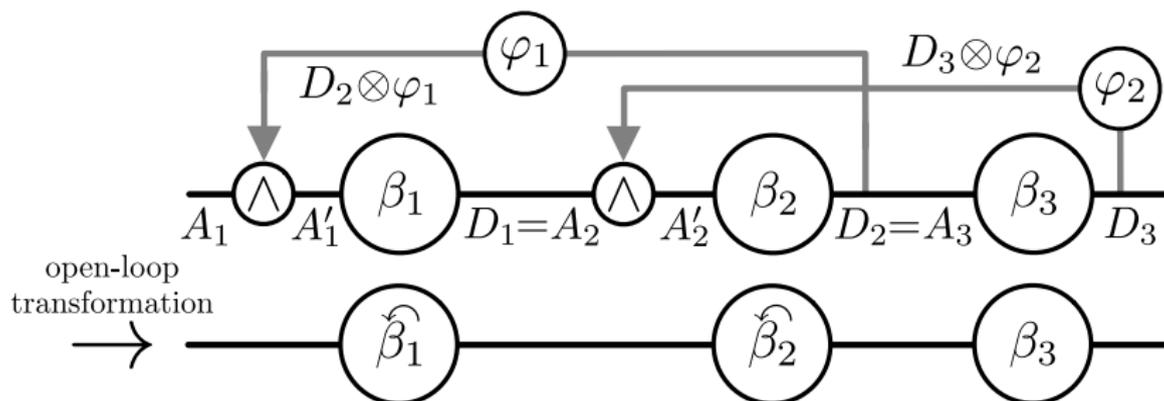


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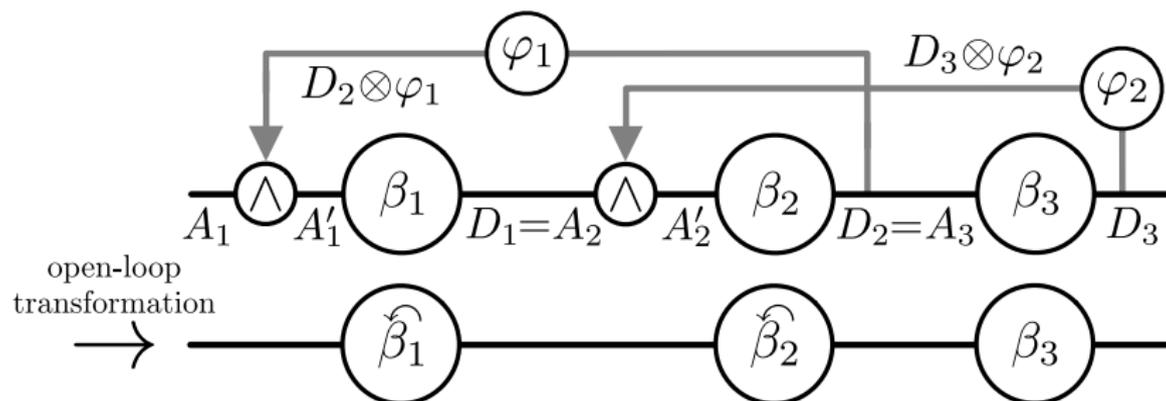


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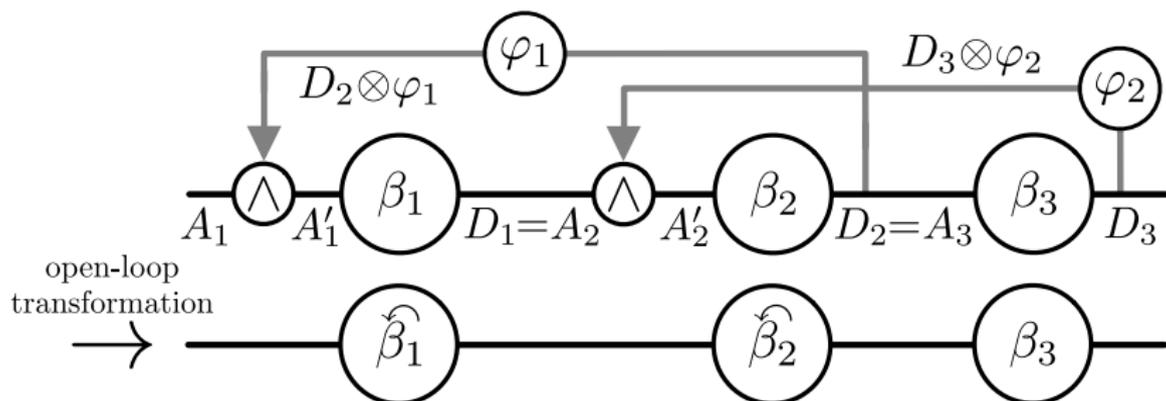


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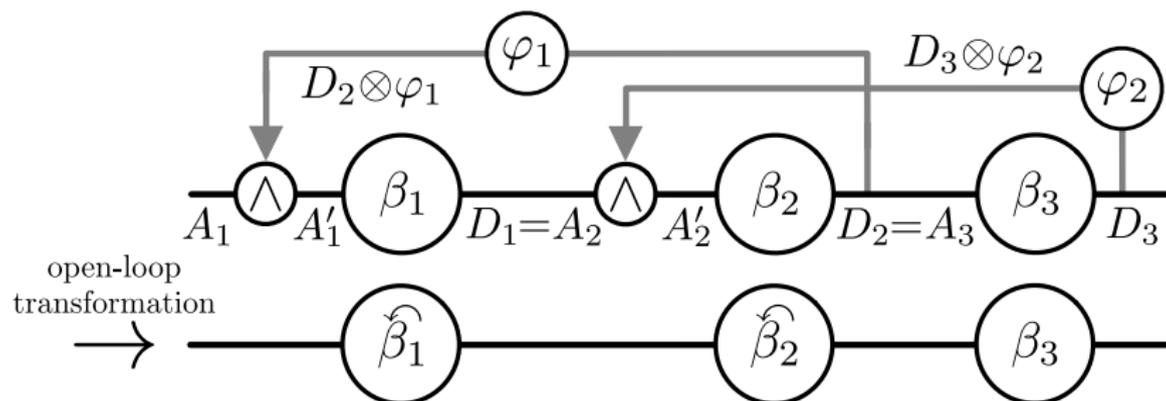


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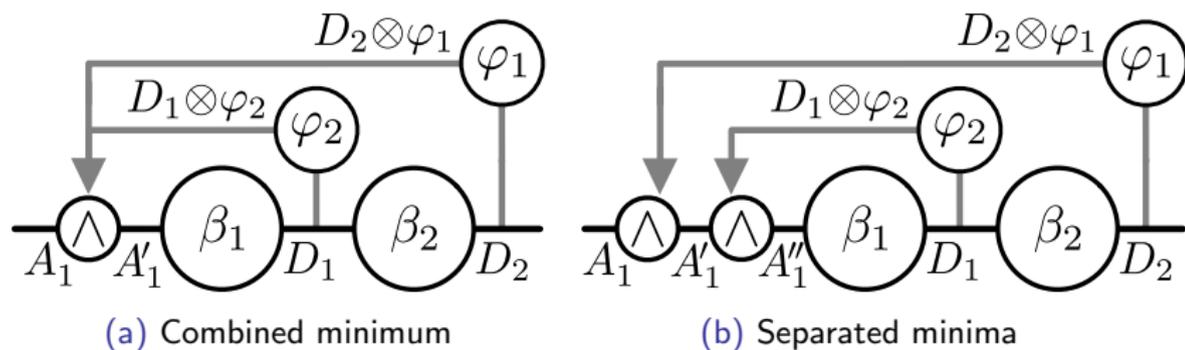


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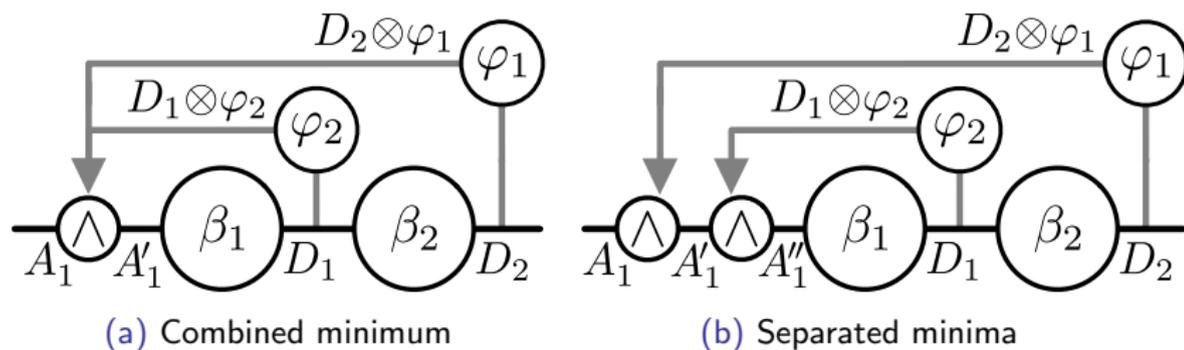


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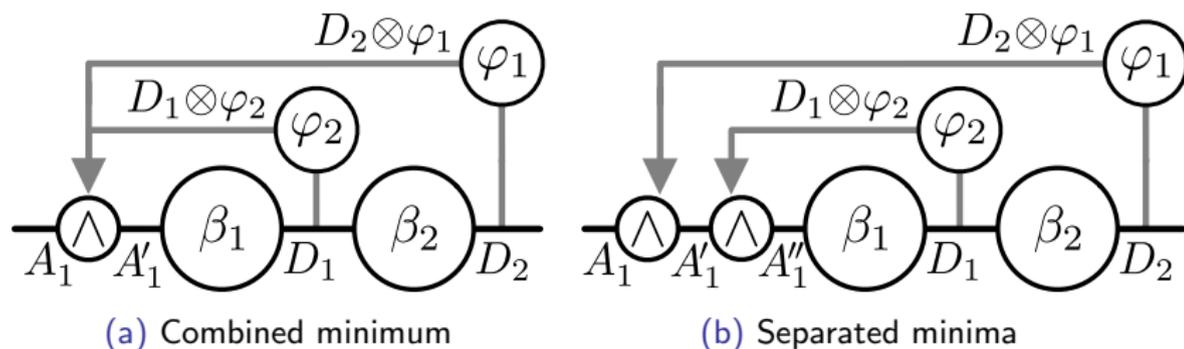


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- Occurs if $i_l = i_m$
- Canonical example: $F_1 = (1, 2, \varphi_1)$, $F_2 = (1, 1, \varphi_2)$
- Can show that structure (a) is equivalent to structure (b)
 - (a) $A'_1 = A_1 \wedge (D_2 \otimes \varphi_1) \wedge (D_1 \otimes \varphi_2)$
 - (b) $A'_1 = A_1 \wedge (D_2 \otimes \varphi_1)$,
 $A''_1 = A'_1 \wedge (D_1 \otimes \varphi_2) = A_1 \wedge (D_2 \otimes \varphi_1) \wedge (D_1 \otimes \varphi_2)$

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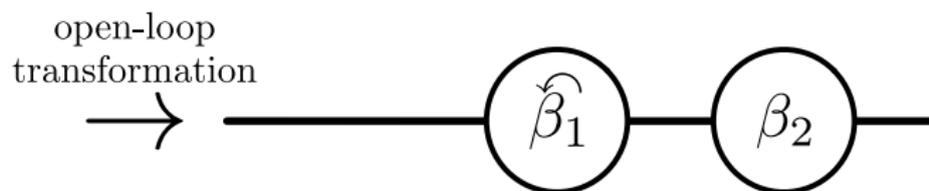


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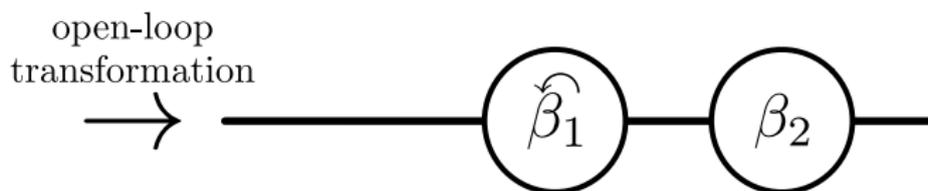


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- Have seen in all canonical example that there is one feedback constraint that can be safely resolved first
- Using the established ordering of feedback constraints, this is always the feedback with the largest ordering index
- General procedure:
 1. Transform feedback constraint F_l with largest index l and exchange β_{i_l} with $\hat{\beta}_{i_l}$
 2. Remove F_l from set of feedback constraints
 3. Repeat until no feedback constraints are left in the system

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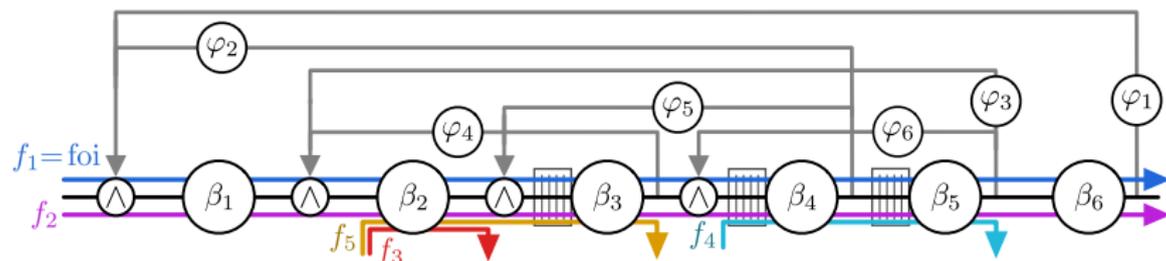


Figure: Example network before transformation.

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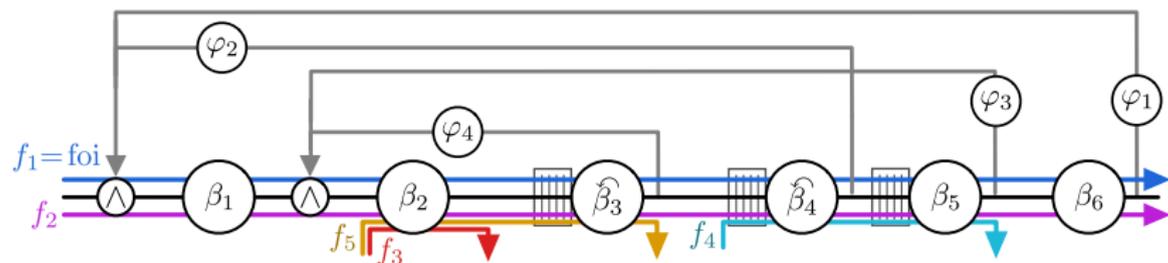


Figure: Second transformation.

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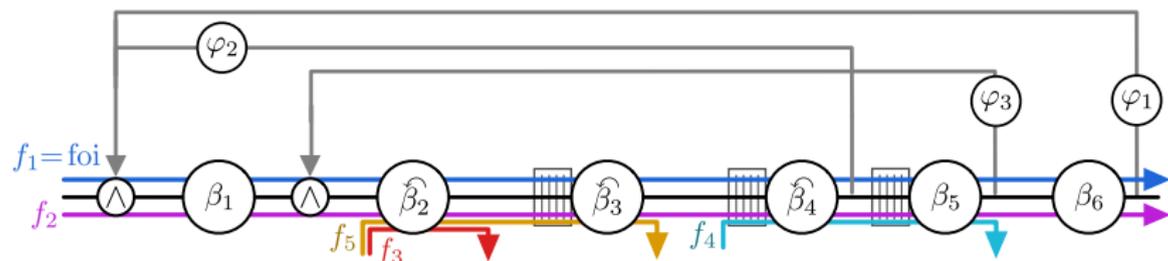


Figure: Third transformation.

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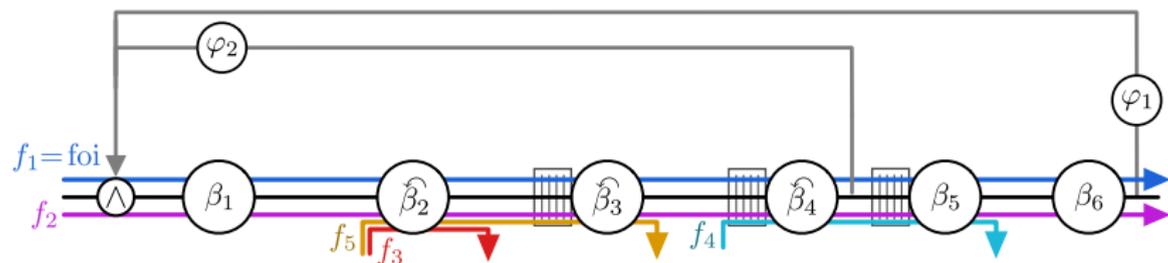


Figure: Fourth transformation.

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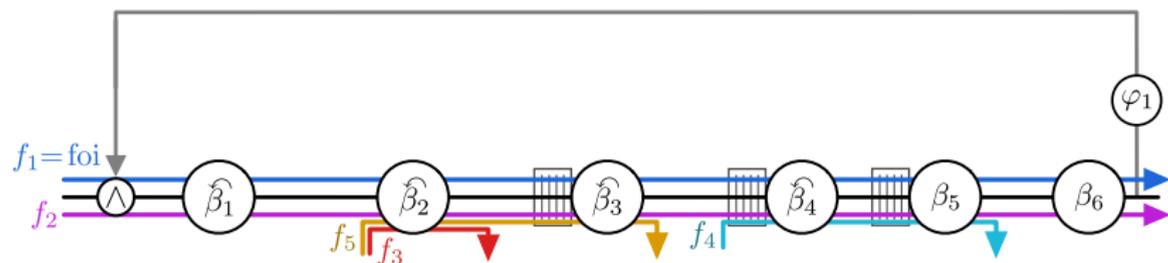


Figure: Fifth transformation.

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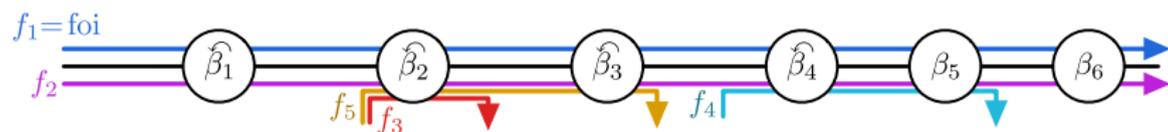


Figure: Transformed open-loop system.

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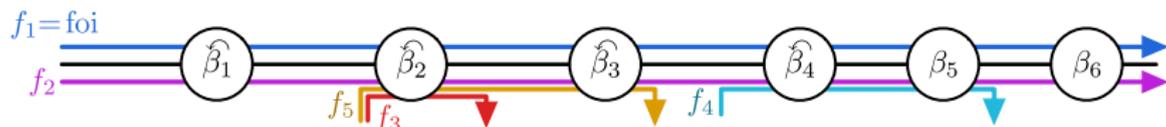


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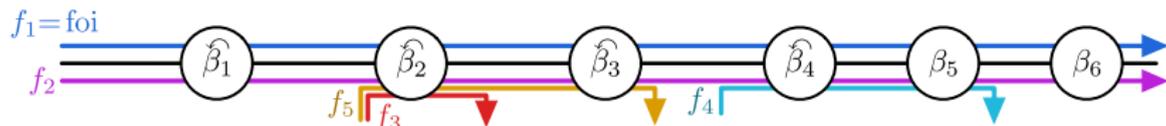


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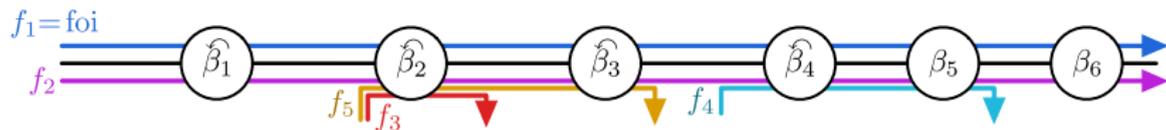


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- Start with an open-loop system, where each transformed node offers a *min-plus* service curve
- Interested in performance bounds for a flow of interest f_1
- Use general residual service curve formula from [Bouillard et al., 2018]

$$\beta_i^{\text{res}} = \beta_i - \bar{\alpha}_i^c$$

with $\bar{\alpha}_i^c$ the maximal arrival curve for the cross flows

Obtaining Per-Flow Performance Bounds

- Use result from [Hamscher et al., 2024] to obtain performance bounds

Theorem

Let an arrival process A traverse a system \mathcal{S} . Further, let the arrivals be constrained by maximal arrival curve $\bar{\alpha} \in \mathcal{F}_0^\uparrow$, and minimal arrival curve $\underline{\alpha} \in \mathcal{F}_0^\uparrow$, and let the system offer a min-plus service curve $\beta \in \mathcal{F}$, with $\xi := \beta_\downarrow$ its lower non-decreasing closure. The backlog $q(t)$ satisfies for all t

$$q(t) \leq v(\bar{\alpha}, \xi) \wedge \sup_{s \geq 0} \{\bar{\alpha}(s)\}.$$

The virtual delay $d(t)$ satisfies for all $t \geq 0$

$$d(t) \leq \max\{z(\underline{\alpha}, \xi), h(\bar{\alpha}, \xi)\}.$$

with $z(\underline{\alpha}, \xi) := \inf\{\tau \geq 0 \mid \underline{\alpha} \otimes \xi(\tau) \geq 0\}$

Conclusion

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 - Idea: build precedence graph for the feedback constraints and find its topological order
- ⇒ Transform feedback constraints in the order provided by the topological order

Closing Remarks

- Identified and showed how different interdependence types between feedback constraints are resolved
 - Here, considered tandem networks
 - Lexicographical order of feedback constraints is enough to identify order of evaluation
 - Order of evaluation for feed-forward networks should be straightforward, but left to be proven
 - Idea: build precedence graph for the feedback constraints and find its topological order
- ⇒ Transform feedback constraints in the order provided by the topological order
- Applying this method to tandem networks provides the correct order of evaluation



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Any questions?

Generalization To Arbitrary Length

- Can generalize the Theorem to transform a system with a feedback constraint of arbitrary length

Proposition

Let a flow f traverse a tandem of n nodes, each offering a service curve $\beta_k, k \in \{1, \dots, n\}$, with a single feedback constraint $F = (i, j, \varphi)$. The tandem offers an end-to-end service curve to the flow f

$$D \geq A \otimes \beta_{e2e}, \quad \text{with } \beta_{e2e} := \bigotimes_{k=1}^{i-1} \beta_k \otimes \hat{\beta}_i \otimes \bigotimes_{l=i+1}^n \beta_l, \quad (1)$$

and a transformed service curve at node i

$$\hat{\beta}_i := \beta_i \otimes \left(\bigotimes_{k=i}^j \beta_k \otimes \varphi \right)^*. \quad (2)$$

Compounded Interdependency

Note that in the case of $i_k = i_l$ and $j_k = j_l$, i.e., the extreme case of compounded feedback spanning the same set of nodes, we only need to evaluate a single feedback loop, however, with feedback constraint $\varphi_k \wedge \varphi_l$. Using distributivity,

$$A' = A \wedge (D \otimes \varphi_k) \wedge (D \otimes \varphi_l) = A \wedge D \otimes (\varphi_k \wedge \varphi_l).$$