Nancy.Expressions A library for automated optimizations

IACOPO CANETTA, ANTONIO A. SALVALAGGIO, ANDREA TRASACCO, <u>RAFFAELE ZIPPO</u>, GIOVANNI STEA UNIVERSITY OF PISA

Raffaele Zippo - Nancy. Expressions - WoNeCa 2025

Nancy.Expressions

A library for computation of Deterministic Network Calculus *expressions*

Design and implementation by Andrea Trasacco for his M.Sc. Thesis

Presented at VALUETOOLS 2024 (proceedings pending) as Nancy. Expressions: Towards a CAS for DNC – A. Trasacco, R. Zippo, G. Stea

Packages, code, documentation already available on Nuget and GitHub



A simple problem, using Nancy

$$(eta_1 - lpha_1)^+_\uparrow \otimes \left(eta_4 - ((lpha_3 + lpha_4) \oslash eta_3))^+_\uparrow
ight)$$
 [BBL18]



var left = Curve.Subtraction(beta1, alpha1)
 .ToNonNegative()
 .ToUpperNonDecreasing();

var right = Curve.Subtraction(beta4, tmp)
 .ToNonNegative()

.ToUpperNonDecreasing();

A simple problem, using Nancy

$$(eta_1 - lpha_1)^+_\uparrow \otimes \left(eta_4 - ((lpha_3 + lpha_4) \oslash eta_3))^+_\uparrow
ight)$$
 [BBL18]

var left = Curve.Subtraction(beta1, alpha1)
 .ToNonNegative()
 .ToUpperNonDecreasing();

var right = Curve.Subtraction(beta4, tmp)
.ToNonNegative()

.ToUpperNonDecreasing();

Some undesirable things:

- Debugging using only the resulting curve, not how are they computed
- Code *is* computation strategy Trade-off between readability/simplicity and clever optimizations
- 3. Some models have *options*, like optimizable *parameters*
 - The above problems scale with program complexity – we usually *compose* studies on networks

The same problem, using Nancy.Expressions

 $(eta_1 - lpha_1)^+_{\uparrow} \otimes (eta_4 - ((lpha_3 + lpha_4) \oslash eta_3))^+_{\uparrow})$ [BBL18]

O(1)Only constructs the *expression tree*



var left = Expressions.Subtraction(beta1, alpha1)
 .ToNonNegative()
 .ToUpperNonDecreasing();

var tmp = Expressions.Addition(alpha3, alpha4)
.Deconvolution(beta3);

var right = Expressions.Subtraction(beta4, tmp)

.ToNonNegative()

.ToUpperNonDecreasing();

1. Debugging using only the resulting curve, not how are they computed

Can print out the expression and its result, separately



2. Code is computation strategy. Trade-off between readability and optimizations

In Nancy.Expressions, most of the program builds the expressions to be computed We can implement many optimizations as <u>transformations</u> or alternative <u>Compute() algorithms</u>

Only the last part of the program needs to be altered to take advantage of these

Examples in the second part of the presentation

3. Some models have options, like optimizable parameters

Expressions can be transformed by replacing sub-expressions

Г

$$\beta_i^{\theta} = \left[\beta - \sum_{j \neq i} \alpha_j \ast \delta_{\theta} \right] \wedge \delta_{\theta}$$

1+

FIFO residual s.c. [BBL18] $\forall \theta \ge 0, \beta_i^{\theta}$ is a s.c.

```
var delta theta 0 = Expressions.FromCurve(new
DelayServiceCurve(0));
```

```
var b_theta_0 = Expressions.Subtraction(
```

beta,

```
Expressions.Addition(contendingFlows)
```

```
.Convolution(delta_theta_0))
```

.ToNonNegative()

```
.Minimum(delta_theta_0);
```

```
var wcd_0 = Expressions.HorizontalDeviation(foi,
b_theta_0).Compute();
```

```
var delta_theta_1 = Expressions.FromCurve(new
DelayServiceCurve(1));
```

```
var wcd_1 = Expressions.HorizontalDeviation(foi,
b_theta_1).Compute();
```

4. The problems scale with program complexity – we compose network studies

We converted a tool using Nancy (from NPBA repo) to Nancy.Expressions. It was straightforward, requiring little understanding of the underlying logic

```
// Using Nancy
public static Curve MakePeriodicWorstCaseArrivalCurve(Flow flow)
//Theorem 1
{
    var linkSpeed = flow.Path[0].LinkSpeed;
    var m = flow.MaxIntervalFrame * flow.MaxFrameSize * 8;
    var linkSpeedCurve = new SigmaRhoArrivalCurve(0, linkSpeed);
    var stair = new StairCurve(m, flow.ClassMeasurementInterval);
    return Curve.Convolution(stair, linkSpeedCurve);
}
```

4. The problems scale with program complexity – we compose network studies

We converted a tool using Nancy (from NPBA repo) to Nancy.Expressions. It was straightforward, requiring little understanding of the underlying logic

hdev(((linkspeedFC-FS ^ th_1brflow0) + (linkspeedFC-FS ^ th_1brflow1) + (linkspeedFC-FS ^ th_1brflow2) + (linkspeedFC-FS ^ th_1brflow3) + (linkspeedFC-FS ^ th_1brflow4) + (linkspeedFC-FS ^ th_1brflow5) + (linkspeedFC-FS ^ th_1brflow6) + (linkspeedFC-FS ^ th_1brflow7) + (linkspeedFC-FS ^ th_1brflow8) + (linkspeedFC-FS ^ th_1brflow9)), minscAFC-FS) + hdev(((shaperAFC-FS ^ ((((linkspeedFC-FS ^ th_1brflow0) + (linkspeedFC-FS ^ th_1brflow1) + (linkspeedFC-FS ^ th_1brflow2) + (linkspeedFC-FS ^ th_1brflow3) + (linkspeedFC-FS ^ th_1brflow4) + (linkspeedFC-FS ^ th_1brflow5) + (linkspeedFC-FS ^ th_1brflow6) + (linkspeedFC-FS ∧ th_1brflow7) + (linkspeedFC-FS ∧ th_1brflow8) + (linkspeedFC-FS ∧ th_1brflow9)) ⊗ maxscAFC-FS) ⊘ minscAFC-FS))), minscAFS-BS) + hdev(((shaperAFS-BS ∧ ((((shaperAFC-FS ∧ ((((linkspeedFC-FS ∧ th_1brflow0) + (linkspeedFC-FS ^ th_1brflow1) + (linkspeedFC-FS ^ th_1brflow2) + (linkspeedFC-FS ^ th_1brflow3) + (linkspeedFC-FS ^ th_1brflow4) + (linkspeedFC-FS ^ th_1brflow5) + (linkspeedFC-FS ^ th_1brflow6) + (linkspeedFC-FS ∧ th_1brflow7) + (linkspeedFC-FS ∧ th_1brflow8) + (linkspeedFC-FS ∧ th_1brflow9)) ⊗ maxscAFC-FS) Ø minscAFC-FS))) ⊗ maxscAFS-BS) Ø minscAFS-BS)) + (shaperAHU-BS ∧ ((((linkspeedHU-BS ∧ th_1brflow10) + (linkspeedHU-BS ^ th_1brflow11) + (linkspeedHU-BS ^ th_1brflow12) + (linkspeedHU-BS ^ th_1brflow13) + (linkspeedHU-BS ^ th_1brflow14) + (linkspeedHU-BS ^ th_1brflow15) + (linkspeedHU-BS ^ th_1brflow16) + (linkspeedHU-BS ∧ th_1brflow17) + (linkspeedHU-BS ∧ th_1brflow18) + (linkspeedHU-BS ∧ th_1brflow19)) ⊗ maxscAHU-BS) ⊘ minscAHU-BS))), minscABS-RU)

Optimizing with Nancy. Expressions

A few use cases and the road ahead

2. Code is computation strategy. Trade-off between readability and optimizations

In Nancy.Expressions, most of the program builds the expressions to be computed We can implement many optimizations as <u>transformations</u> or alternative <u>Compute() algorithms</u> Only the last part of the program needs to be altered to take advantage of these

Equivalences

Computation Scheduling

Alternative Algorithms

Transform expressions for speed of computation



Take the computation path that takes the least



Different *Compute()* algorithms to get the same result



Optimizing with Equivalences

Many results in literature simplify expressions into easier-to-compute ones Instead of doing it with pen-and-paper, Nancy.Expressions allows to automate it

Example from literature [BJLL06, ZS23], involving $O(n^2)$ convolutions

$$\beta^{eq'} = \bigotimes_{i=1}^{n-1} \left(\beta_i \bigotimes_{j=i}^{n-1} (\overline{\beta_j \otimes \beta_{j+1} + W_{j+1}}) \right) \otimes \beta_n$$

The convolution \otimes is commutative and associative, and $\overline{f} \otimes \overline{f} = \overline{f}$ So in [BJLL06, ZS23], this was manually optimized into O(n) convolutions

$$\beta^{eq'} = \left(\bigotimes_{i=1\dots n} \beta_i\right) \otimes \bigotimes_{i=1\dots n-1} \overline{\beta_i \otimes \beta_{i+1} + W_{i+1}}.$$

Optimizing with Equivalences

Equivalence objects represent these properties – including their hypotheses

We can use them iteratively to simplify expressions

```
do
{
    previous = expression;
    expression = previous.ApplyEquivalence(equivalence);
} while (expression != previous);
```

Scales well in instances where optimizations are hard to spot or apply

Optimizing with Equivalences

We benchmarked a simple program implementing the intuitive $O(n^2)$ expression



Fig. 4: Average execution time and memory allocation of the end-to-end service curve of a flow-controlled network with n nodes.

Optimizing with computation scheduling

This result was explored by Antonio A. Salvalaggio in his B.Sc. thesis

Presented as Improving (min,+) Convolutions by Heuristic Operation Reordering at JWRTC @RTNS24, won JRWRTC Best Paper Award

- 1. Associative and commutative operations yield the same result regardless of order of computation
- 2. But the *computation time* may not be the same

This work focused on (min,+) convolution and considered only its commutativity



Optimizing with computation scheduling

We studied convolutions of 6 operands, comparing the computation time of all permutations



Optimizing with computation scheduling

We experimented with combining heuristics for *first-pair* and *next-curve* selection

The resulting algorithm can pick an order of computations that is "to the left" of the distribution

Better than random ordering, with little overhead



Current work is aiming to generalize this approach to consider associativity as well The key are the *heuristics* used to predict cost and explore the options

Optimizing with *Compute()* algorithms

Before the call to Compute(), the program builds the expression tree object with all the necessary data

We can use this to use different strategies – like **Finitary RTC**

Finitary RTC was discussed in a series of papers [GY13, LBS16, LBSGY17] Can be applied to *hdev()* and *vdev()* expressions, under some hypotheses Generalized as an expression-based algorithm by **lacopo Canetta** in his M.Sc. thesis Main problem: period explosion in large studies



Finitary RTC was discussed in a series of papers [GY13, LBS16, LBSGY17] Can be applied to hdev() and vdev() expressions, under some hypotheses Generalized as an expression-based algorithm by **lacopo Canetta** in his M.Sc. thesis Main problem: period explosion in large studies

Core idea: bound the domain of interest, removing the unused periodic parts





Finitary RTC was discussed in a series of papers [GY13, LBS16, LBSGY17] Can be applied to hdev() and vdev() expressions, under some hypotheses Generalized as an expression-based algorithm by **lacopo Canetta** in his M.Sc. thesis Main problem: period explosion in large studies

Core idea: bound the domain of interest, removing the unused periodic parts



A 4-step algorithm:

1. Replace each operand with UA approximations



Top-down phase

A 4-step algorithm:

- **1**. Replace each operand with UA approximations
- 2. Compute bound on hdev say K_{hdev}



A 4-step algorithm:

- **1**. Replace each operand with UA approximations
- 2. Compute bound on hdev say K_{hdev}
- 3. Propagate to find per-operand bounds say K_{β_1} , K_{α_1} , ...
 - Replace each operand with its finite cut, e.g. $\alpha'_1 = \alpha_1([0, K_{\alpha_1}])$



Top-down phase

A 4-step algorithm:

- **1**. Replace each operand with UA approximations
- 2. Compute bound on hdev say K_{hdev}
- 3. Propagate to find per-operand bounds say K_{β_1} , K_{α_1} , ...
 - Replace each operand with its finite cut, e.g. $\alpha'_1 = \alpha_1([0, K_{\alpha_1}])$
- 4. Run the computation on finite cuts
 - Same result, no period explosion



Exploiting this is problematic: you need to rewrite a great deal of code

var c2AlphaU Gurw = Gpc.ComputeAlphaU(c1AlphaU, beta2CutU, beta2CutL); var c2BetaU Gurw = Gpc.ComputeBetaU(c1AlphaL, beta2CutU); var c2BetaL Gurw = Gpc.ComputeBetaL(c1AlphaU, beta2CutL); Console.WirteLine("Compute G2");

var c3BetaL Curve = Gpc.ComputeBetaL(c2AlphaU, beta3CutL); Console.WriteLine("Computed C3");

var c5AlphaU Conve = Gpc.ComputeAlphaU(c4AlphaU, c2BetaU); var c5BetaU Conve = Gpc.ComputeBetaU(c4AlphaL, c2BetaU); var c5BetaL Conve = Gpc.ComputeBetaL(c4AlphaU, c2BetaL); Console.WirteLine("Compute G5");

var c6BetaL :Curve = Gpc.ComputeBetaL(c5AlphaU, c3BetaL); Console.WriteLine("Computed C6");

var c8AtphaU_cSeve = 6pc.ComputeAtphaU(c7AtphaU, c5BetaU, c5BetaU); var c8BetaU_CSeve = 6pc.ComputeBetaU(c7AtphaL, c5BetaU); var c8BetaU_Seve = 6pc.ComputeBetaL(c7AtphaU, c5BetaL); Console.WiteLine("ComputeD 68");

var c9BetaL Curve = Gpc.ComputeBetaL(c8AlphaU, c6BetaL); Console.WriteLine("Computed C9");

var c10AlphaU Curve = Gpc.ComputeAlphaU(alpha4CutU, c7BetaU, c7BetaL); Console.WriteLine("Computed C10");

var c11AlphaU Curve = 6pc.ComputeAlphaU(c10AlphaU, c8BetaU, c8BetaL); Console.WriteLine("Computed C11");

var delay10 Rational = Curve.HorizontalDeviation(alpha4CutU, c78etaL); var delay11 Rational = Curve.HorizontalDeviation(c18AlphaU, c88etaL); var delay12 Rational = Curve.HorizontalDeviation(c11AlphaU, c98etaL);

Nancy, without F-RTC, 35 lines of code

var c3BetaLinU = Gpc.ComputeBetaU(c2AlphaLinL, beta3LinU); var c3BetaLinL = Gpc.ComputeBetaL(c2AlphaLinU, beta3LinL); Console.WriteLine("Computed L C3");

var c6BetaLinU = Gpc.ComputeBetaU(c5AlphaLinL, c3BetaLinU); var c6BetaLinL = Gpc.ComputeBetaL(c5AlphaLinU, c3BetaLinL); Console.WriteLine("Computed L C6");

var (c7AlphaLinU, c7AlphaLinL,c7BetaLinU, c7BetaLinL) =
 Gpc.ComputerGpc(alpha3LinU, alpha3LinL, c4BetaLinU, c4BetaLinL);
 Console.WriteLine("Computed L (7");

var c9BetaLinU = Gpc.ComputeBetaU(c8AlphaLinL, c6BetaLinU); var c9BetaLinL = Gpc.ComputeBetaL(c8AlphaLinU, c6BetaLinL); Console.WriteLine("Computed L C9");

var c10AlphaLinU = Gpc.ComputeAlphaU(alpha4LinU, c78etaLinU, c78etaLinL); var c10AlphaLinL = Gpc.ComputeAlphaL(alpha4LinL, c78etaLinU, c78etaLinL);

Console.WriteLine("Computed L C10");

var c11AlphaLinU = 6pc.ComputeAlphaU(c18AlphaLinU, c8BetaLinU, c8BetaLinU); var c11AlphaLinL = 6pc.ComputeAlphaL(c18AlphaLinL, c8BetaLinU, c8BetaLinU); Console.WriteLine("Computed L C11");

///doi/of all 8 (respective) all 8 (respective)

kaC12, kaC12, kaC12, kaC10, kaC100, kaC100, law (); var (xC104)pha1m, kC10041a1m, kC10041a1m, kC10041a1m) = dpc,FintTaryBoundaries (alpha41im, a)pha41im, (Pha41im, CPHe41im, A) Retonal.Res(VaC11, kC114)pha100(), RetOnal.Res(kaC11, kC154)pha1n(), RetOnal.Res(KaC11, kC114)pha100(), RetOnal.Res(kaC11, kC154)pha10(), RetOnal.Res(KaC11, kC114)pha10(), RetOnal.Res(KaC11, kC154)pha10(), RetOnal.Res(KaC11, kC114)pha10(), RetOnal.Res(KaC114)pha10(), RetOnal.Res(KaC114)pha10(), RetOnal.Res(KaC114)pha10(), RetOnal.Res(KaC114)pha10(), RetOnal.Res(KaC114)pha10(), RetOnal.Res(KaC114)pha10(), RetOnal.Res(KaC114)pha10(), RetOnal.Res(KaC114)pha10(), RetOnal.Res(KaC114)pha10(

var (kC9AlphaInU, kC9AlphaInL, kC9BetaInU, kC9BetaInL) = Gpc.FinitaryBoundar (c6AlphaInU, kC9AlphaInL, kC9BetaInU, kC9BetaInL, Rational.Jene, Rational.Jene, kC12, kC(2):

Rational.Zero, Rational.Zero, KXC12, KXC12); var (KCBAlphaInU, KCBAlphaInU, KCBBetaInU, KCBBetaInU) = Gpc.FinitaryBoundaries

(c7Alphalinu, c7Alphalinu, c5Metalinu, C5Metalinu, K50Alphalinu, K50Alphalinu, K50Alphalinu, Kational.Max(bbCll, K11Betalinu), Rational.Max(bbCll, KC1Betalinu)); var (K57Alphalinu, K57Alphalinu, K57Betalinu, K5Metalinu) - Gpc.FinitaryBoundaries

- (ar) Activation and Activation (arbitrario) and a second arbitrary boundaries (alphalinu, alphalini, celetalini, description, arbitrario), Arbitrario, Arbitrario, KC00etalo, KC00etalo, KC00etalo, KC00etalo, KC00etalo, KC00etalo), KC00etalo, KC00etalo
- var (kCGAlphaInU, KCGAlphaInL, KCGBetaInU, KCGBetaInL) Gpc.FinitaryBoundaries (cSAlphaLinU, cSAlphaLinL, cIBetaLinU, cIBetaLinL,
- (cfAlphaLinu, cfAlphaLinu, cHBetaLinu, cHBetaLinu, Rational.Zero, Rational.Zero, kC9BetaInU, kC9BetaInL);

- var (kcSAlphaInU, kcSAlphaInL, kcSBetaInU, kcSBetaInL) = 6pc.Finitary8oundaries (c4AlphaInU, c4AlphaInL, c2BetaInU, c2BetaInL, kcGAlphaInU, kc6AlphaInL, kcBBetaInU, kc8BetaInL);
- var (kCAAlphaInU, kCAAlphaInL, kCABetaInU, kCABetaInL) = Gpc.FinitaryBoundaries (alpha2LinU, alpha2LinL, cIBetaInU, cIBetaInL, kCSAlphaInU, kCSAlphaInL, kC7BetaInU, kC7BetaInL);
- var (kC3AlphaInU, kC3AlphaInL, kC3BetaInU, kC3BetaInL) = Gpc.FinitaryBoundaries (c2AlphaInU, c2AlphaInL, beta3LinU, beta3LinL, Rational.Zero, Rational.Zero, Kc6BetaInU, kC6BetaInL);
- var (kC2AlphaInU, kC2AlphaInL, kC2BetaInU, kC2BetaInL) = Gpc.Finitary8oundaries (cIAlphaLinU, cIAlphaLinL, betaZLinU, betaZLinL, kC3AlphaInU, kC3AlphaInL, kC5BetaInU, kC5BetaInL);
- var (kClAlphaInU, kClAlphaInL, kClBetaInU, KClBetaInL) = Gpc.FinitaryBoundaries (alphaILinU, alphaILinL, betaILinU, betaILinL, kC2AlphaInU, kC2AlphaInL, kC4BetaInU, kC4BetaInL);
- var alpha1CutU = alpha1.UpperCurve.FinitaryCurve(kC1AlphaInU); var alpha2CutU = alpha2.UpperCurve.FinitaryCurve(kC4AlphaInU);
- var alpha3CutU = alpha3.UpperCurve.FinitaryCurve(kC7AlphaInU); var alpha1CutL = alpha1.LowerCurve.FinitaryCurve(kC1AlphaInL);
- var alpha2CutL = alpha2.LowerCurve.FinitaryCurve(kC4AlphaInL);
- var alpha3CutL = alpha3.LowerCurve.FinitaryCurve(kC7AlphaInL);

var k4U = Rational.Max(kaC10, kC10AlphaInU); var alpha4CutU = alpha4.UpperCurve.FinitaryCurve(k4U);

van betaiutU = betal.UpperCurve.FinitaryCurve(kCiBetaInU); van beta2utU = betal.UpperCurve.FinitaryCurve(kCiBetaInU); van betaIutL = betal.LowerCurve.FinitaryCurve(kCiBetaInL); van beta2utL = betal.LowerCurve.FinitaryCurve(kCiBetaInL); van beta2utL = betal.uperCurve.FinitaryCurve(kCiBetaInL);

var c2AlphaU = Gpc.ComputeAlphaU(c1AlphaU, beta2CutU, beta2CutL); var c2BetaU = Gpc.ComputeBetaU(c1AlphaL, beta2CutU); var c2BetaL = Gpc.ComputeBetaL(c1AlphaU, beta2CutL); Console.MriteLine("Computed C2");

var c3BetaL = Gpc.ComputeBetaL(c2AlphaU, beta3CutL); Console.WriteLine("Computed C3");

var (c4AlphaU, c4AlphaL,c4BetaU, c4BetaL) =
 Gpc.ComputerGpc(alpha2CutU, alpha2CutL, c1BetaU, c1BetaL);
Console.WriteLine("Computed C4");

var c5AlphaU = Gpc.ComputeAlphaU(c4AlphaU, c2BetaU, c2BetaL); var c5BetaU = Gpc.ComputeBetaU(c4AlphaL, c2BetaU); var c5BetaL = Gpc.ComputeBetaL(c4AlphaU, c2BetaL); Console.WriteLine("Computed C5");

var c6BetaL = Gpc.ComputeBetaL(c5AlphaU, c3BetaL); Console.WriteLine("Computed C6");

var (c7AlphaU, c7AlphaL,c7BetaU, c7BetaL) =
 Gpc.ComputerGpc(alpha3CutU, alpha3CutL, c4BetaU, c4BetaL);
Console.WriteLine("Computed C7");

var c8AlphaU = Gpc.ComputeAlphaU(c7AlphaU, c5BetaU, c5BetaL); var c8BetaU = Gpc.ComputeBetaU(c7AlphaL, c5BetaU); var c8BetaL = Gpc.ComputeBetaL(c7AlphaL, c5BetaL); Console.WriteLine("Computed C8");

var c9BetaL = Gpc.ComputeBetaL(c8AlphaU, c6BetaL); Console.WriteLine("Computed C9");

var c10AlphaU = Gpc.ComputeAlphaU(alpha4CutU, c7BetaU, c7BetaL); Console.WriteLine("Computed C10");

var c11AlphaU = Gpc.ComputeAlphaU(c10AlphaU, c8BetaU, c8BetaL); Console.WriteLine("Computed C11");

var delay10 = Curve.HorizontalDeviation(alpha4CutU, c78etal); var delay11 = Curve.HorizontalDeviation(18Alpha4), c88etal); var delay12 = Curve.HorizontalDeviation(c11AlphaU, c98etaL); var delay = delay10 + delay11 + delay12; Console.Hriteline(\$"The delay is: (delay)");

Nancy, with F-RTC, 139 lines of code

Raffaele Zippo - Nancy. Expressions - WoNeCa 2025

With the expression-based algorithm and Nancy. Expressions, this is now free

var result = delayExpression.Compute();

var result = delayExpression.ComputeFinitary();

var (clAlphaU CurreExpression , clAtphaL CurreExpression , clBetaU CurreExpression , clBetaL CurreExpression) = EppEExpressionComponent.6pcCurreExpression(alphalU, alphalL, betalU, betalL); Console.WriteLine("Computed Cl");

var c2AlphaU CurreEuression = GpcExpressionComponent.AlphaUExpression(clAlphaU, beta2U); var c2BetaU CurreEuression = GpcExpressionComponent.BetaUExpression(clAlphaU, beta2U); var c2BetaU CurreEuression = GpcExpressionComponent.BetaUExpression(clAlphaU, beta2U); Console.WriteLine("Computed C2");

var c3BetaL CurveExpression = GpcExpressionComponent.BetaLExpression(c2AlphaU, beta3L); Console.WriteLine("Computed C3");

var (c4AlphaU CurveExpression , c4AlphaL CurveExpression , c4BetaU CurveExpression , c4BetaL CurveExpression) = EppEExpressionComponent.EppECurveExpression(alpha2U, alpha2L, c1BetaU, c1BetaL); Console.HirtLine("Computed 64");

var cSALphaU ConvExpession = GpcExpressionComponent.AtphaUExpression(c4ALphaU, c2BetaU); var cSBetaU ConveExpession = GpcExpressionComponent.BetaUExpression(c4ALphaL, c2BetaU); var cSBetaU ConveExpession = GpcExpressionComponent.BetaUExpression(c4ALphaU, c2BetaU); Comosle.WriteLine('Computed C5');

var c68etaL :CurveExpression = GpcExpressionComponent.BetaLExpression(c5AlphaU, c3BetaL); Console.WriteLine("Computed C6");

var (c7AlphaU CurreExpression , c7AlphaL CurreExpression , c7BetaU CurreExpression , c7BetaL CurreExpression) = EppEExpressionComponent.EppEurveExpression(alpha3U, alpha3L, c4BetaU, c4BetaL); Console.FirtLine("Computed C7");

var c8ALphaU CorreEpression = GpcExpressionComponent.ALphaUExpression(c7ALphaU, c5BetaU); var c8BetaU CorreEpression = GpcExpressionComponent.BetaUExpression(c7ALphaL, c5BetaU); var c8BetaL CorreEpression = GpcExpressionComponent.BetaLExpression(c7ALphaU, c5BetaU); Console.MriteLine('Computed (SP');

var c9BetaL=CurveExpression = GpcExpressionComponent.BetaLExpression(c8AlphaU, c6BetaL); Console.WriteLine("Computed C9");

var c10AlphaU.CurveExpression = GpcExpressionComponent.AlphaUExpression(alpha4U, c7BetaU, c7BetaL); Console.WriteLine("Computed C10");

var c11AlphaU/CurveExpression = GpcExpressionComponent.AlphaUExpression(c10AlphaU, c8BetaU, c8BetaL); Console.WriteLine("Computed C11");

var delay10 = new HorizontalDeviationExpression(alpha4U, RightExpression c78etaL); var delay11 = new HorizontalDeviationExpression(c18AlphaU, RightExpression(c88etaL);

var detay11 = new HorizontalDeviationExpression(cloktpnad, HighExpression coBetaL); var detay12 = new HorizontalDeviationExpression(cl1AlphaU, RighExpression c9BetaL);

var delay Rational = delay10.ComputeFinitary() + delay11.ComputeFinitary() + delay12.ComputeFinitary();

Nancy.Expressions, 35 lines of code

Now that it is easy to use – is it always worth it?

No.



time

Now that it is easy to use – is it always worth it?

No. The finite prefix from F-RTC may be larger than the *minimal representation*



Now that it is easy to use – is it always worth it?

No. The finite prefix from F-RTC may be larger than the *minimal representation*

Increasing the burst or the rate utilizations are easy ways to find examples where plain UPP curves are much faster – using Nancy optimizations like super-isospeed

Next goal: find heuristic to transparently apply the best method



Nancy.Expressions

A library for computation of Deterministic Network Calculus *expressions*

Open-source, written in C#, extends the existing *Nancy* library

While Nancy implements operators, Nancy. Expressions generalizes it to expressions

- Build, visualize and manipulate expressions before computing them
- Apply equivalences to simplify expressions
- Use different computation algorithms without rewriting the entire program







For Teaching

For Research

For Applications

Nancy.Expressions

Developer experience

Writing and debugging analysis tool is improved thanks to the focus on expressions, that remove a lot of noise from the program structure

State of the art accessibility

Results from literature like theorems, properties and compute algorithms can be coded in and re-used much more easily

New paths for optimizations

Our experiments with the above showed many trade-off instances

There is much room to research heuristics to improve runtime and make harder or more accurate studies feasible

Thanks!



References

[BJLL06] A. Bose, X. Jiang, B. Liu, and G. Li, *Analysis of Manufacturing Blocking Systems with Network Calculus*, Performance Evaluation, 2006

[GY13] N. Guan and W. Yi, *Finitary Real-Time Calculus: Efficient Performance Analysis of Distributed Embedded Systems*, IEEE RTSS 2013

[LBS16] K. Lampka, S. Bondorf and J. B. Schmitt, *Achieving Efficiency without Sacrificing Model Accuracy: Network Calculus on Compact Domains*, IEEE MASCOTS 2016

[LBSGY17] K. Lampka, S. Bondorf, J. B. Schmitt, N. Guan and W. Yi, *Generalized Finitary Real-Time Calculus*, IEEE INFOCOM 2017

[ZS22] R. Zippo and G. Stea, Nancy: an efficient parallel Network Calculus library, SoftwareX, 2022

[ZS23] R. Zippo and G. Stea, *Computationally Efficient Worst-Case Analysis of Flow-Controlled Networks With Network Calculus*, IEEE Transactions on Information Theory, 2023

[SZS24] A. A. Salvalaggio, R. Zippo, and G. Stea, *Improving (min,+) Convolutions by Heuristic Operation Reordering*, JWRTC @RTNS24

[TZS24] A. Trasacco, R. Zippo, and G. Stea, Nancy. Expressions: Towards a CAS for DNC, VALUETOOLS 2024

Equivalences and *cut-through* properties

An important part of applying an equivalence is <u>testing its hypotheses</u> Example: knowing that f is subadditive, we can simplify $f \otimes f = f$ To test subadditivity, we need to *verify* that $f \otimes f = f$ Recall that f may be an *expression*

$$\beta^{eq'} = \bigotimes_{i=1}^{n-1} \left(\beta_i \bigotimes_{j=i}^{n-1} (\overline{\beta_j \otimes \beta_{j+1} + W_{j+1}}) \right) \otimes \beta_n$$
$$\beta^{eq'} = \left(\bigotimes_{i=1\dots n} \beta_i \right) \otimes \bigotimes_{i=1\dots n-1} \overline{\beta_i \otimes \beta_{i+1} + W_{i+1}}.$$

Equivalences and *cut-through* properties

Core idea: while testing properties, many computations can be skipped if the results in literature are cleverly used



The convolution is (surely) subadditive if its operands are The subadditive closure is subadditive The sum is (surely) subadditive if its operands are Token-bucket arrival curves are subadditive

So we can conclude this expression is subadditive without computing it and its self-convolution

Nancy.Expressions is designed to take advantage of these inferences avoiding unnecessary computations

Optimizing with Equivalences - Syntax

In code

Using an ad-hoc syntax

```
var equivalence = new Equivalence(
                                                                             {
    leftSideExpression: Expressions.Convolution(
                                                                                 f in U subadditive;
        Expressions.Placeholder("f"),
                                                                                 g in U zero-at-zero;
        Expressions.Placeholder("g")),
    rightSideExpression: Expressions.Placeholder("f")
                                                                                 f <= g;
);
                                                                                 f * g well-defined;
equivalence.AddHypothesis("f", f => f.IsSubAdditive);
                                                                             } == > f * g = f
equivalence.AddHypothesis("g", g => g.IsZeroAtZero);
equivalence.AddHypothesis(
    "f", "g", (CurveExpression f, CurveExpression g) => f <= g);
equivalence.AddHypothesis(
```

"f", "g", (f, g) => f.Convolution(g).IsWellDefined);

Optimizing with Equivalences - Limitations

In its current state, application is automated but **manually triggered**.

The user is expected to know which equivalences may help

A valuable goal of this is to *transparently* apply known equivalences to optimize This is a non-trivial, with different challenges

- What is the cost of looking for matches and verifying the hypotheses?
- What is the time saved by using the equivalence?
- What about substitutions that enable *other* equivalences?
- How should we limit the exploration overhead vs. the computation time?

Optimizing with computation scheduling - insights

The difference comes from wasted computations

Nancy can minimize the result *after* a computation, but also uses optimizations that skip unnecessary ones *before*

A "good path" has good intermediate operations that avoid useless computations that are later discarded

Addressing associativity is more complicated, but has several advantages

- A larger search space should include lower minima
- Can exploit parallelization on multicore systems
- Can be mixed with other results, e.g. convolutions of convex/concave/subadditive functions

An abstraction hierarchy

Abstraction Level	Possible optimizations	Software
DNC and system model	How the system model is expressed into curves, e.g. a strict service curve can be replaced by its superadditive closure.	Your analysis tool, ??
(min,+) and (max,+) expressions	Use algebraic properties to compute the <i>expression</i> faster	Nancy.Expressions, NC Maude?
(min,+) and (max,+) operations	Use algebraic properties to compute the <i>operation</i> faster	Nancy, RTC toolbox